



4.4.3 Backus-Naur Form

Backus-Naur Form

Backus-Naur form, like regular expressions, is another way of representing languages.

Unlike regular expressions, **BNFs can make use of recursion** which allows it to represent more languages.

Production Rules

BNF can use production rules to describe a language.

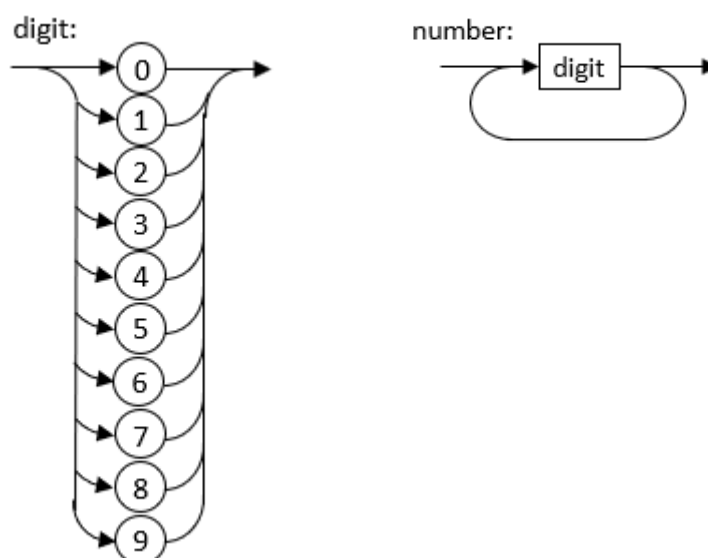
$$\langle \text{Digit} \rangle ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$$
$$\langle \text{Number} \rangle ::= \langle \text{Digit} \rangle \mid \langle \text{Digit} \rangle \langle \text{Number} \rangle$$

In the example above:

- The definition of a Digit uses $\mid$ as an 'or' and is defined as any value from 0 to 9.
- The definition of a Number uses recursion to build numbers with one or more digits

Syntax Diagrams

Syntax diagrams give a visual representation of production rules.



A circle represents a terminal symbol.

A rectangle represents a non-terminal.

The digit non-terminal is used when defining a number.





BNF recursion

Using recursion BNF can describe languages that could not be described using regular expressions.

An example of this would be a language that has an 'a' enclosed between any number of opening and closing brackets.

$$\{ a, (a), ((a)), (((a))), (((((a))))), \dots \}$$

There is no way of representing this language using a regular expression but using recursion in Backus-Naur form:

$$\langle \text{Term} \rangle ::= a \mid (\langle \text{Term} \rangle)$$

Caution

While using recursion allows Backus-Naur form to define languages that a regular expression cannot, just because a production rule is recursive doesn't mean that a regular expression could not define the same language.

$$\langle \text{Symbol} \rangle ::= a \mid b$$

$$\langle \text{Word} \rangle ::= \langle \text{Symbol} \rangle \mid \langle \text{Symbol} \rangle \langle \text{Word} \rangle$$

The production for Word is recursive but describes the same language as the regular expression $(a|b)^+$.

